Ad Hoc Wireless Networks Key Management for Heterogeneous

Chandar Sekhar Das, Ashok Muduli, Satyabrata Das, Rasmi Sarangi

Department of Computer science and Engineering, NM Institute of Engineering and Technology, Bhubaneswar, Odisha

Department of Computer science and Engineering, Raajdhani Engineering College,Bhubaneswar,Odisha Department of Computer science and Engineering ,Aryan Institute of Engineering and Technology Bhubnaeswar, Odisha

Department of Computer science and Engineering, Capital Engineering College

I. INTRODUCTION

Since its birth more than two decades ago, public key cryptography has been recognized as one of the most ef- fective mechanisms for providing fundamental security ser- vices including authentication, digital signatures and en- cryption for dynamic networks. Effective management of keys, or digital certificates holding the keys, is one of the key factors for the successful wide-spread deployment of public key cryptography. PKI (Public Key Infrastructure), an infrastructure for managing digital certificates, was in- troduced for this purpose. The most important component of PKI is the CA (Certificate Authority), the trusted entity in the system that vouches for the validity of digital certifi- cates. The success of PKI depends on the availability of the CA to the principals in the system (or the nodes in the net- work) since a principal must correspond with the CA to get a certificate, check the status of another principal's certifi- cate, acquire another principal's digital certificate, and so on. PKI has been widely deployed for wired networks and some infrastructure-based wireless networks. Since good connectivity can be assumed in such networks, the main thrust of research in such environments has focused on the availability of the CA and the scalability of the CA to handle a large number of requests.

However, it is still unclear if such approaches can be ex- tended to ad hoc networks. Connectivity, which was as- sumed to be good in previous PKI solutions, is no longer stable in ad hoc networks. Unfortunately, maintaining con- nectivity is one of the main challenges in ad hoc networks, since the inherent infrastructurelessness of ad hoc networks makes it hard to guarantee any kind of connectivity. An- other serious problem present in ad hoc networks is the physical vulnerability of the nodes themselves. Consider- ing that most ad hoc networks will be deployed with mobile nodes, the possibility of the nodes being captured or compromised is higher than in wired networks with stationary

This work was supported in part by the National Science Foundation under grant ANI-0081308. hosts. Mobile nodes in a infrastructure-based wireless net- works have the same vulnerability, but they can rely on the infrastructure for detection of compromised nodes and po- tentially some help with recovery. With an infrastructure- based solution, mobile nodes may store all sensitive infor- mation in the infrastructure and maintain minimal informa- tion in the device. Since there is no stable entity in an ad hoc network, the vulnerability of ad hoc nodes is even higher.

Currently proposed solutions for providing PKI for ad hoc networks address the physical vulnerability of the nodes by employing the distribution of CA functionality across multiple nodes and using threshold cryptography[4, 2]. These approaches also increase the availability of the CA. The first challenge in such approaches is picking which nodes should be a part of the distributed CA. A random set of nodes from a network can be chosen to play the role of CA. Such a random choice may not be best for security of the whole network. The second challenge is how to provide efficient, yet effective, communication between the mobile nodes and the CA nodes to create the illusion of an available CA, even in dynamic networks with possible compromises or partitions.

This work describes a framework to provide efficient yet effective distributed CA service for ad hoc wireless net- works. We select physically or computationally more se- cure nodes as MOCAs (MObile Certificate Authority) and use threshold cryptography[1] to distribute the CA's private key among these MOCA nodes. We also provide a protocol for clients to contact MOCAs and get certification services without incurring excessive overhead.

II. MOCA APPROACH

The main technique in our approach is to use threshold cryptography to distribute the CA's private key among many nodes that collectively act as the CA for the network. Given this distributed approach, there are two main challenges: (1) How to pick the MOCA nodes and (2) How to make them available for services.

In our work, MOCAs are chosen based on an observation of heterogeneity within an ad hoc network. Most existing work blindly considers all nodes in an ad hoc network to be identical, which is not necessarily true.

For example, in a military battle field scenario (one of MANET's most pop- ular examples), there can be many different types of mo- bile nodes in the field (e.g. infantry soldiers, tanks, platoon leader's jeeps, command and control vehicles, etc). This heterogeneity can be exploited when choosing the MOCA nodes. Physically more secure and computationally more powerful nodes are the typical choices for MOCAs. These selected MOCAs share the secret and collectively provide CA functionality.

Client nodes are equipped with MP(MOCA certification Protocol), which enables communication with a sufficient number of MOCAs in an efficient and effective way. In initial version of MP, clients flood the network with certi- fication requests and MOCAs that receive the request re- ply with a certification reply. This approach works very well in terms of success ratio but has high overhead due to its flooding nature. To alleviate this overhead, we in-vestigated the cache tables of client nodes and discovered that with a certain amount of certification traffic in the net- work, a mobile node tends to have many cached routes en- tries to enough MOCA nodes. To exploit this, we developed a unicast-based optimization of MP that uses multiple sets of unicasts without route discovery when there are enough cached routes in the client's local cache. Since there can be more than enough cached routes, we tested three different selection criteria: closest MOCAs, freshest MOCAs, and random MOCAs.

III. PRELIMINARY RESULTS

We use NS-2 network simulator for our simulations. Our initial flooding based approach shows a good success ratio but suffers from overhead due to flooding.

Fig. 1 shows the number of successfully received certifica- tion replies for 1000 certification requests. 30 MOCAs are deployed in the network, so the maximum number of replies for a certification request is 30. Under varying mobility pat- terns, more than 90% percent of requests are answered with 15 or more replies.

While all three of our unicast-based protocols work more efficiently than flooding, closest-unicast that picks the clos- est MOCAs from the client turns out to be the best in suc- cess ratio and latency. Fig. 2 shows the comparison of suc- cess ratios between different flavors of unicast-based pro- tocols. Flooding is shown as the baseline. More detailed simulation results are available in [3].

Our simulation results show the effectiveness of our ap- proach and we provide some insights into the configuration of such security services in ad hoc networks.

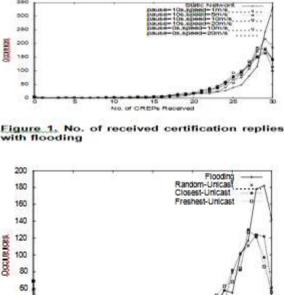


Figure 2. No. of received certification replies with unicast-base protocols

15

No. of CREPs Received

10

20

25

30

40 20 0

0

5

REFERENCES

- [1] Y. Frankel and Y. G. Desmedt. Parallel Reliable Thresh- old Multisignature. Technical Report TR-92-04-02, Univ. of Wisconsin–Milwaukee, 1992.
- [2] J. Kong, P. Zerfos, H. Luo, S. Lu, and L. Zhang. Providing Robust and Ubiquitous Security Support for Mobile Ad-Hoc Networks. In Proceedings of ICNP '01.
- [3] S. Yi and R. Kravets. Key Management for Heterogeneous Ad Hoc Wireless Networks. Technical Report UIUCDCS-R- 2002-2290/UILU-ENG-2002-1734, University of Illinois at Urbana-Champaign, 2002.
- [4] L. Zhou and Z. J. Haas. Securing Ad Hoc Networks. IEEENetwork Magazine, Nov. 1999.